

Lecture 8 - Sep 30

Math Review

***Rel Image vs. Func Application
Modelling: Rel vs. Partial vs. Total Func
Injection, Surjection, Bijection***

Announcements/Reminders

- Today's class: [notes template](#) posted
- **Event-B Summary** Document
- Priorities:
 - + **Lab1** → Review
 - + **Lab2** → Review
- Change of **ProgTest** venue – WSC106/108
- Released:
 - + **ProgTest** guide
 - + 2 Practice Tests and solutions
 - + **Lab1**, **Lab2** solutions

Relational **Image** vs. Functional **Application**

A function is a **relation**.

$$f \in \boxed{S} \rightarrow \boxed{T}$$

$$f = \{ (3, a), (1, b) \}$$

relation
image

Exercises:

$$f[\{3\}] = \{a\}$$

$$f[\{1\}] = \{b\}$$

$$f[\{2\}] = \emptyset$$

$\in S \notin \text{dom}(f)$

functional application

$$f(3) = a$$

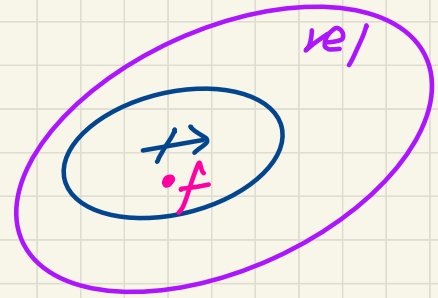
$$f(1) = b$$

$$f(2) = \text{undefined}$$

$$\text{isFunction}(f) \wedge \text{dom}(f) \subseteq \{1, 2, 3\}$$

$$\text{dom}(f) \subseteq \{1, 2, 3\}$$

at least one value from S that does not have the corresponding range value.



In Rodin

$$\textcircled{1} f : \mathbb{Z} \leftrightarrow \mathbb{Z}$$

$$f[\{c\}]$$

$$c \in \mathbb{Z}$$

always well-defined

$$\textcircled{2} g : \mathbb{Z} \rightarrow \mathbb{Z}$$

$$g(c)$$

$$c \in \mathbb{Z}$$

not always well-defined
 $\hookrightarrow$ one PO generated:
 $c \in \text{dom}(g)$.

Cardinality of
relational image

$$r \in S \leftrightarrow T$$

r is also a function

$$\Rightarrow |r[\{s\}]| \leq 1$$

$\downarrow$
 $s \in S$

$\hookrightarrow \begin{cases} = 1 \leadsto s \in \text{dom}(r) \\ = 0 \leadsto s \notin \text{dom}(r) \end{cases}$

Modelling Decision: Relations vs. Functions

An organization has a system for keeping track of its employees as to where they are on the premises (e.g., ``Zone A, Floor 23``). To achieve this, each employee is issued with an active badge which, when scanned, synchronizes their current positions to a central database.

Assume the following two sets:

- *Employee* denotes the **set** of all employees working for the organization.
- *Location* denotes the **set** of all valid locations in the organization.

Is $\text{where_is} \in \text{Employee} \leftrightarrow \text{Location}$ appropriate?

No. $e_1 \mapsto l_1 \in \text{where_is} \wedge e_1 \mapsto l_2 \in \text{where_is}$

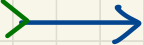
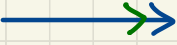
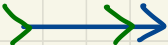
Is $\text{where_is} \in \text{Employee} \rightarrow \text{Location}$ appropriate?

↳ No $\because$ some employees may not be in the company.

Is $\text{where_is} \in \text{Employee} \rightarrow \text{Location}$ appropriate?

↓
YES.

Functions

	(dom) injective	(ran) surjective	(dom, ran) bijective
partial →	partial injection 	partial surjection 	n.a.
total →			

Contrapositive: $s_1 \neq s_2 \Rightarrow \neg((s_1, t) \in f \wedge (s_2, t) \in f)$

Injective Functions

func. property
 $isInjective(f)$

$$\forall s, t_1, t_2 \bullet (s \in S \wedge t_1 \in T \wedge t_2 \in T) \Rightarrow ((s, t_1) \in f \wedge (s, t_2) \in f \Rightarrow t_1 = t_2)$$

$$\forall s_1, s_2, t \bullet (s_1 \in S \wedge s_2 \in S \wedge t \in T) \Rightarrow ((s_1, t) \in f \wedge (s_2, t) \in f \Rightarrow s_1 = s_2)$$

same range value

inj. property
 If f is a **partial injection**, we write: $f \in S \rightsquigarrow T$

- e.g., $\{\emptyset, \{(1, a)\}, \{(2, a), (3, b)\}\} \subseteq \{1, 2, 3\} \rightsquigarrow \{a, b\}$
- e.g., $\{(1, b), (2, a), (3, b)\} \notin \{1, 2, 3\} \rightsquigarrow \{a, b\}$
- e.g., $\{(1, b), (3, b)\} \notin \{1, 2, 3\} \rightsquigarrow \{a, b\}$

If f is a **total injection**, we write: $f \in S \rightarrow T$

- e.g., $\{1, 2, 3\} \rightarrow \{a, b\} = \emptyset$
- e.g., $\{(2, d), (1, a), (3, c)\} \in \{1, 2, 3\} \rightarrow \{a, b, c, d\}$
- e.g., $\{(2, d), (1, c)\} \notin \{1, 2, 3\} \rightarrow \{a, b, c, d\}$
- e.g., $\{(2, d), (1, c), (3, d)\} \notin \{1, 2, 3\} \rightarrow \{a, b, c, d\}$

cannot have two distinct dom. values mapping to the same range value.

e.g. $\{(1, b), (2, b)\}$
 $\downarrow \quad \quad \downarrow$
 $s_1 \quad t \quad s_2 \quad t$
 $\hookrightarrow ((1, b) \in f) \wedge ((2, b) \in f) \Rightarrow 1=2$
false

func^(V) prop. $\wedge$ injective^(V) prop.

If f is a **partial injection**, we write: $f \in S \twoheadrightarrow T$

- e.g., $\{\emptyset, \{(1, a)\}, \{(2, a), (3, b)\}\} \subseteq \{1, 2, 3\} \twoheadrightarrow \{a, b\}$
- e.g., $\{(1, b), (2, a), (3, b)\} \notin \{1, 2, 3\} \twoheadrightarrow \{a, b\}$ * func ✓
- e.g., $\{(1, b), (3, b)\} \notin \{1, 2, 3\} \twoheadrightarrow \{a, b\}$ X

inj X $\because$ distinct dom values
1 and 3 both
map to b.

If f is a **total injection**, we write: $f \in S \rightarrow T$

- e.g., $\{1, 2, 3\} \rightarrow \{a, b\} = \emptyset$
- e.g., $\{(2, d), (1, a), (3, c)\} \in \{1, 2, 3\} \rightarrow \{a, b, c, d\}$
- e.g., $\{(2, d), (1, c)\} \notin \{1, 2, 3\} \rightarrow \{a, b, c, d\}$
- e.g., $\{(2, d), (1, c), (3, d)\} \notin \{1, 2, 3\} \rightarrow \{a, b, c, d\}$

sat. both
func. & inj.
properties (V)
initially

distinct dom values
map to distinct
ran values $\Rightarrow$ injective.

$$\{(\underline{1}, \underline{a}), (\underline{2}, \underline{a})\}$$

↳ function
not injective

$$\{(1, a), (1, b)\}$$

↳ not a function!

$S \longleftrightarrow T$: set of all relations $\rightarrow$ set of all possible total injections.

If f is a **total injection**, we write: $f \in S \rightarrow T$

- e.g., $\{1, 2, 3\} \rightarrow \{a, b\} = \emptyset$
- e.g., $\{(2, d), (1, a), (3, c)\} \in \{1, 2, 3\} \rightarrow \{a, b, c, d\}$
- e.g., $\{(2, d), (1, c)\} \notin \{1, 2, 3\} \rightarrow \{a, b, c, d\}$
- e.g., $\{(2, d), (1, c), (3, d)\} \notin \{1, 2, 3\} \rightarrow \{a, b, c, d\}$

$\{(1, \underline{a}), (2, \underline{b}), (3, \underline{a})\}$
 $\downarrow$ violates inj. prop.

WITNESS

	func. prop	total	inj. prop
①	✓	✓	✓
②	✓	✗	✓
③	✓	✓	✗

$$f \in S \leftrightarrow T$$

Surjective Functions

$$\text{total}(f) \iff \text{dom}(f) = S$$

$$\text{isSurjective}(f) \iff \text{ran}(f) = T$$

① funct pop.
② surj prop.

If f is a **partial surjection**, we write: $f \in S \twoheadrightarrow T$

- e.g., $\{ \{(1, b), (2, a)\}, \{(1, b), (2, a), (3, b)\} \} \subseteq \{1, 2, 3\} \twoheadrightarrow \{a, b\}$
- e.g., $\{(2, a), (1, a), (3, a)\} \notin \{1, 2, 3\} \twoheadrightarrow \{a, b\}$ $\rightarrow$ func not surj.
- e.g., $\{(2, b), (1, b)\} \notin \{1, 2, 3\} \twoheadrightarrow \{a, b\}$

If f is a **total surjection**, we write: $f \in S \rightarrow T$

- e.g., $\{ \{(2, a), (1, b), (3, a)\}, \{(2, b), (1, a), (3, b)\} \} \subseteq \{1, 2, 3\} \rightarrow \{a, b\}$
- e.g., $\{(2, a), (3, b)\} \notin \{1, 2, 3\} \rightarrow \{a, b\}$
- e.g., $\{(2, a), (3, a), (1, a)\} \notin \{1, 2, 3\} \rightarrow \{a, b\}$

	total	func	surj.	
①	✓	✓	✓	✓
②	✓	✓	✓	✓
③	x	✓	✓	x
④	✓	✓	x	x

Bijjective Functions

f is **bijjective/a bijection/one-to-one correspondence** if f is **total**, **injective**, and **surjective**.

all possible bijections

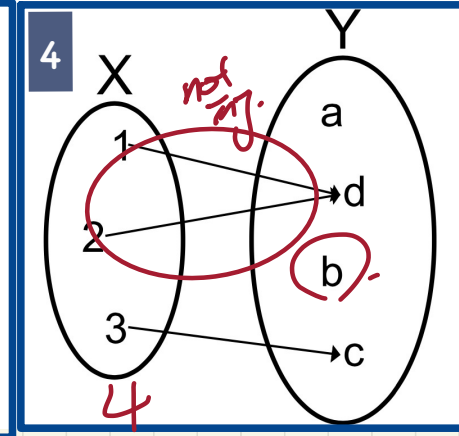
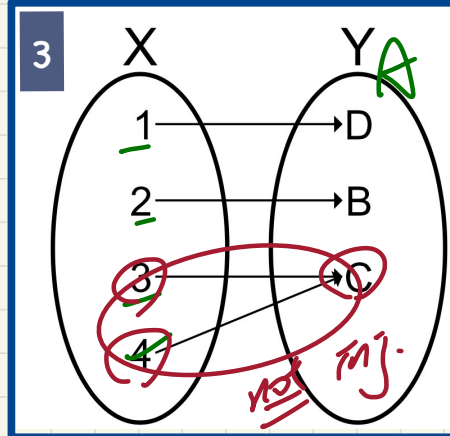
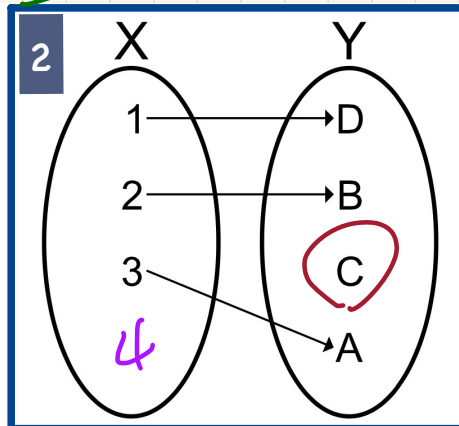
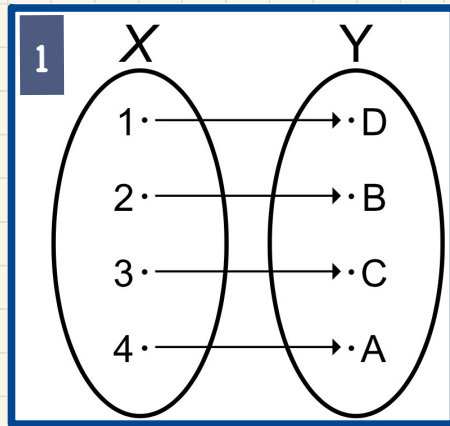
- e.g., $\{1, 2, 3\} \rightsquigarrow \{a, b\} = \emptyset \rightsquigarrow \because$ no injective function can be made
- e.g., $\{ \{(1, a), (2, b), (3, c)\}, \{(2, a), (3, b), (1, c)\} \} \subseteq \{1, 2, 3\} \rightsquigarrow \{a, b, c\}$
- e.g., ① $\{(2, b), (3, c), (4, a)\} \notin \{1, 2, 3, 4\} \rightsquigarrow \{a, b, c\}$
- e.g., ② $\{(1, a), (2, b), (3, c), (4, a)\} \notin \{1, 2, 3, 4\} \rightsquigarrow \{a, b, c\}$
- e.g., ③ $\{(\underline{1}, \underline{a}), (\underline{2}, \underline{c})\} \notin \{\underline{1}, \underline{2}\} \rightsquigarrow \{\underline{a}, \underline{b}, \underline{c}\}$

	total + func	inj	surj.	
①	X	✓	✓	X
②	✓	X	✓	X
③	✓	✓	X	X

Exercise

$$X = \{1, 2, 3, 4\}$$

$$Y = \{A, B, C, D\}$$



	1	2	3	4
partial	✓	✓	✓	✓
total	✓	X	✓	X
injection	✓	✓	X	X
surjection	✓	X	X	X
bijection	✓	X	X	X